## Math 325 Topics Sheet for Exam 1

Throughout, we rely on set notation to make our ideas precise. Sets are typically described as the collection of all objects from a specific 'universe' that meet certain specific conditions: E.g.,  $S = \{x \in \mathbb{Q} \mid x^2 < 2\}$  = the rational numbers (universe!) whose square is less than 2 (condition!).

 $A \subseteq B$  means that every element of A is also an element of B; A = B is typically established by showing that  $A \subseteq B$  and  $B \subseteq A$ .

 $A \cup B = \{x \mid x \in A \text{ textrmor } x \in B\} \text{ ; } A \cap B = \{x \mid x \in A \text{ textrmand } x \in B\} \text{ ; } A^c = \{x \mid x \notin A\}$ 

Functions: a function  $f : A \to B$  is a rule which assigns to each  $x \in A$  (the *domain*) exactly one element  $f(x) \in B$  (the *codomain*).

We will mostly focus on functions with  $A \subseteq \mathbb{R}$  and  $B \subseteq \mathbb{R}$ . Functions come in different flavors:

increasing: if x > y then f(x) > f(y). non-decreasing:  $x \ge y$  implies  $f(x) \ge f(y)$ .

decreasing: if x > y then f(x) < f(y). non-increasing:  $x \ge y$  implies  $f(x) \le f(y)$ .

One-to-one: if f(x) = f(y), then x = y. Alternate form: if  $x \neq y$  then  $f(x) \neq f(y)$ . Third form!: for any  $y \in B$ , there is at most one  $x \in A$  with f(x) = y.

Onto: for every  $y \in B$ , there is <u>at least</u> one  $x \in A$  with f(x) = y.

One-to-one <u>and</u> onto = a one-to-one correspondence (or *bijection*).

Note: one-to-one and onto have a lot to do with what the domain and codomain of the funciton f are!

Composition:  $f: A \to B$ ,  $g: B \to C$ , then  $g \circ f: A \to C$  is  $(g \circ f)(x) = g(f(x))$ 

If  $g \circ f$  is onto, then g is onto! If  $g \circ f$  is one-to-one, then f is one-to-one!

The Real Line: Nearly everything we will do comes down to understanding the properties of the real line  $\mathbb{R}$ .

 $\mathbb{R} = a$  complete ordered field

**Field:** we have addition + and multiplication  $\cdot$ , so that

(A1) addition exists: if  $x, y \in \mathbb{R}$  then  $x + y \in \mathbb{R}$ 

(A2) commutativity: x + y = y + x for every  $x, y \in \mathbb{R}$ 

(A3) associativity: x + (y + z) = (x + y) + z for every  $x, y, z \in \mathbb{R}$ 

(A4) zero exists: there is  $0 \in \mathbb{R}$  so that x + 0 = x for every  $x \in \mathbb{R}$ 

(A5) additive inverses exist: for every  $x \in \mathbb{R}$  there is a  $(-x) \in \mathbb{R}$  with x + (-x) = 0

(M1) multiplication exists: if  $x, y \in \mathbb{R}$  then  $x \cdot y \in \mathbb{R}$ 

(M2) commutativity:  $x \cdot y = y \cdot x$  for every  $x, y \in \mathbb{R}$ 

(M3) associativity:  $x \cdot (y \cdot z) = (x \cdot y) \cdot z$  for every  $x, y, z \in \mathbb{R}$ 

(M4) one exists: there is  $1 \in \mathbb{R}$  so that  $x \cdot 1 = x$  for every  $x \in \mathbb{R}$ 

(M5) multiplicative inverses exist:

for every 
$$x \in \mathbb{R}$$
 with  $x \neq 0$  there is a  $x^{-1} \in \mathbb{R}$  with  $x \cdot x^{-1} = 1$   
(D) distributivity: for every  $x, y, z \in \mathbb{R}$  we have  $x(y+z) = xy + xz$ 

**Ordered:** there is a collection  $\mathcal{P}$  = the *positive numbers* so that

- (P1) closed under addition: if  $x, y \in \mathcal{P}$  then  $x + y \in \mathcal{P}$
- (P2) closed under multiplication: if  $x, y \in \mathcal{P}$  then  $x \cdot y \in \mathcal{P}$
- (P3) trichotomy:

for every  $x \in \mathbb{R}$  exactly one of the following is true:  $x \in \mathcal{P}$ , or  $-x \in \mathcal{P}$ , or x = 0.

We then defined the ordering x < y to mean that  $y - x \in \mathcal{P}$ . (So x > 0 means  $x \in \mathcal{P}$ .) Basic properties:

trichotomy! for any  $x, y \in \mathbb{R}$  exactly one of x < y, x = y, or x > y is true.

transitivity: x < y and y < z implies that x < z

x < y implies that x + z < y + z for any  $z \in \mathbb{R}$ 

x < y and z > 0 implies that xz < yz (because yz - xz = (y - x)z with  $y - x, z \in \mathcal{P}$ ) weak inequality:  $a \leq b$  means a < b or a = b; similar properties hold!

From these basic properties we can recover many familiar properties we have seen before; for example

(-x)y = -(xy) (-x)(-y) = xy -(-x) = x  $(x^{-1})^{-1} = x$ x < 0 and y > 0 implies xy < 0 , z < 0 and w < 0 implies zw > 0the additive inverse of a number is unique (i.e., if x + y = 0 = x + z then y = z)

**Proving things:** Our ultimate goal is to provide <u>proofs</u> of some of the important results from calculus. This means that we need to <u>justify</u> the assertions we make, showing how a hypothesis forces our collusions to be true. Two often-used approaches:

Case analysis: Starting from a hypothesis (e.g.,  $x \neq 0$ ), one of several possibilities (cases) must be true (e.g., x > 0 or x < 0). If we show that in each case our hoped-for conclusion is true (e.g.,  $x^2 > 0$ ), then the hypothesis implies the conclusion ( $x \neq 0$  implies  $x^2 > 0$ ).

Proof by contradiction: "A implies B" is the same as "it is not possible for A to be true and also that B is false". Proof by contradiction consists of starting from 'A is true and B is false' and showing that we must inevitably show that something we know is false is <u>true</u>. This means that we cannot have A true and B false; so A implies B !

Example: using the Rational Roots Theorem (see below) we can show that it is not possible to have  $x^3 = 5$  and  $x \in \mathbb{Q}$ . So  $x^3 = 5$  implies  $x \notin \mathbb{Q}$ .

Another approach we will often use: induction! (see below)

**Completeness:** From the natural (= counting) numbers  $\mathbb{N}$  we get the integers  $\mathbb{Z}$  (by taking additive inverses) and then the rationals  $\mathbb{Q}$  (by taking multiplicative inverses). But to get the reals  $\mathbb{R}$  we need to step beyond the properties above.

A set  $A \subseteq \mathbb{R}$  is bounded (bdd) from above if there is a  $M \in \mathbb{R}$  so that  $x \leq M$  for every  $x \in A$ .

A least upper bound  $\lambda$  is an upper bound for A so that no smaller number is an upper bound. In symbols:  $x \leq \lambda$  for every  $x \in A$  and if  $\mu < \lambda$  then there is an  $x \in A$  with  $\mu < x$ [Equivalently:  $\lambda$  is an upper bd for A and if  $\nu$  is also an upper bound for A then  $\lambda \leq \nu$ .]

Completeness Axiom: Every non-empty set  $A \subseteq \mathbb{R}$  that is bdd from above has a least upper bound.

Least upper bound of A is unique!  $\lambda = \sup(A)$ 

Application: If  $x, y \in \mathbb{R}$  and y - x > 1, then there is an  $n \in \mathbb{Z}$  with  $x \le n < y$ .

Application:  $A = \{x \in \mathbb{R} : x^2 < 2\}$  is non-empty and bdd above:  $\lambda = \sup(A)$ . Then we showed:  $\lambda^2 = 2$  (!) So  $\lambda =$  what we would call  $\sqrt{2}$ 

Rational Roots Theorem: If  $p(x) = a_0 x^n + \cdots + a_{n-1} x + a_n$  is a polynomial with integer coefficients  $a_i \in \mathbb{Z}$  for all i), and if  $r = \alpha/\beta$  is a rational root of  $p(p(\alpha/\beta) = 0$  where  $\alpha$  and  $\beta$  have no factors in common, then  $\alpha$  evenly divides  $a_n$  and  $\beta$  evenly divides  $a_0$ .

Since  $\sqrt{2}$  is a root of  $p(x) = x^2 - 2$ , which by the rational roots theorem has <u>no</u> rational roots,  $\sqrt{2} \notin \mathbb{Q}$ . (!) By the same reasoning, if  $n \in \mathbb{N}$  and  $\sqrt{n} \in \mathbb{Q}$  then  $\sqrt{n} \in \mathbb{N}$ .

Well-ordering Property: If  $A \subseteq \mathbb{N}$  is non-empty, then it has a *smallest element*. That is, there is an  $\lambda \in A$  so that  $lambda \leq z$  for every  $z \in A$ . (In general for a set A whose infimum  $\lambda$  lies in A, we call it the *mminimum* of A. If the supremum lies in A, we call it the *maximum*.)

Application: the Principle of Mathematical Induction: If  $A \subseteq \mathbb{N}$  is a set satisfying (1)  $n_0 \in A$ , and (2) if  $n \ge n_0$  and  $n \in A$  then  $n + 1 \in A$ , then  $\{n \in \mathbb{N} : n \ge n_0\} \subseteq A$ .

[Why? If not all  $n \ge n_0$  are in A, then there is a smallest n which is not in A. But then  $n-1 \in A$  (or  $n = n_0$ ) both of which contradict our hypotheses!]

PMI as it is usually stated: If P(n) is a statement about the integer n so that

(1)  $P(n_0)$  is <u>true</u>, and

(2) if  $n \ge n_0$  and P(n) is true then P(n+1) must also be true,

<u>then</u> P(n) is true for every integer  $n \ge n_0$ .

Sample applications:

For every integer 
$$n \ge 1$$
 we have  $\sum_{k=1}^{n} k = \frac{n(n+1)}{2}$   
For every integer  $n \ge 1$  we have  $\sum_{k=1}^{n} k^2 = \frac{n(n+1)(2n+1)}{6}$ 

For every  $n \ge 5$  we have  $2^n > n^2$ 

Fibonacci sequence:  $a_0 = a_1 = 1$  and for  $n \ge 1$  we have  $a_{n+1} = a_n + a_{n-1}$ . Then for every  $n \ge$ (something) we have  $a_n \ge \left(\frac{3}{2}\right)^n$ , and for every  $n \ge 12$  we have  $a_n \ge n^2$ .

[For all of these the <u>method</u> by which we established the inductive step is probably more important as we go forward than the actual result!]

More Well-ordered property consequences:

If  $\epsilon > 0$  then there is an  $n \in \mathbb{N}$  so that  $n\epsilon > 0$  (the Archimedean Principle).

[Add a positive number to itself enough times and you will get a big number!] Rationals are everywhere: If  $x, y \in \mathbb{R}$  with x < y then there is an  $r \in \mathbb{Q}$  so that x < r < y.

## Sequences and convergence: Start with *distance*:

Absolute value: |x| = x if  $x \ge 0$ , otherwise it is -x. So  $|x| \ge 0$  for all x.

"Triangle inequality":  $|x + y| \le |x| + |y|$  for every  $x, y \in \mathbb{R}$ .

[Useful 'opposite' consquence:  $|x - y| \ge |x| - |y|$ 

(useful for showing that |x - y| is <u>not</u> small!]

|x - y| = the distance between x and y. Triangle inequality:  $|x - z| \le |x - y| + |y - z|$ Sequences: A sequence is essentially just a function  $f : \mathbb{N} \to \mathbb{R}$ . We write  $f(n) = a_n$ Examples: the Fibonnaci sequence!  $a_0 = a_1 = 1, a_2 = 2, a_3 = 3, a_4 = 5, a_5 = 8, a_6 = 13$ 

$$a_n = 1 + (-1)^n / n$$
,  $b_n = \frac{n^2 - n + 13}{3n^2 + 2n - 13}$   
 $a_n = \sum_{k=1}^n \frac{1}{k^2}$ 

Convergence: What happens to a sequence as n gets large? E.g., for  $b_n = \frac{n^2 - n + 13}{3n^2 + 2n - 13}$ , observation/intuition/calculus suggests that for large  $n \ b_n$  is approximately 1/3. Formalizing this notion leads us to:

 $\lim_{n \to \infty} a_n = L \text{ means that } a_n \text{ is } \underline{\text{close to }} L \text{ so long as } n \text{ is } \underline{\text{large}}.$  More precisely,  $a_n$  is as close as we want it to be to L, so long as n is  $\underline{\text{large enough}}.$  This leads us to

 $\lim_{n \to \infty} a_n = L \text{ means for every } \epsilon > 0 \text{ [think: small!] there is an } N \in \mathbb{N} \text{ [}N \text{ depends on } \epsilon \text{ (!)}\text{]}$ so that  $n \ge N$  implies that  $|a_n - L| < \epsilon$ . If this limit exists (i.e., there is an L that the sequence converges to) we say that the sequence is *convergent*. If there is no number that the sequence converges to, we say that the sequence is *divergent*. [Shorthand: we write  $a_n \to L$ .]

The key here is that we wish to show that  $|a_n - L|$  is small; we typically do this by comparing  $|a_n - L|$  to other things that we <u>know</u> are small. This, in turn, we usually do by altering  $|a_n - L|$ , making it larger (but not too large!), until we end up with something we can show is small (so long as n is large enough!). To help us do this we have some basic limit results:

If  $a_n \to L$  and  $b_b \to M$ , and  $k \in \mathbb{R}$  is a konstant, then:  $ka_n \to kL$   $a_n + b_n \to L + M$   $a_n - b_n \to L - M$   $a_nb_n \to LM$  $a_n/b_n \to L/M$  (so long as  $b_n, M \neq 0$ )

Other properties we learned along the way:

If  $(a_n)_{n=1}^{\infty}$  is convergent, then the set  $A = \{a_n : n \in \mathbb{N}\}$  is bounded. Limits are unique! If  $a_n \to L$  and  $a_n \to M$ , then L = M.

$$\frac{1}{n} \to 0$$

If  $a_n \ge a$  for every n (eventually!) and  $a_n \to L$ , then  $L \ge a$ If  $a_n \ge b_n$  for every n and  $a_n \to L$ ,  $b_n \to M$ , then  $L \ge M$ If  $a_n \to L$  and  $L \ne 0$ , then eventually  $|a_n| > \frac{|L|}{2}$ 

Showing convergence without knowing the limit: monotonicity

 $a_n$  is monotone increasing if  $a_{n+1} \ge a_n$  for every n

 $b_n$  is monotone decreasing if  $b_{n+1} \leq b_n$  for every n

If it is one or the other, we say the sequence is *monotone* (or *monotonic*).

A monotone increasing sequence that is bounded above (i.e., for some  $M \in \mathbb{R}$  we have  $a_n \leq M$  for every  $n \in \mathbb{N}$ ) converges it limit is  $\sup\{a_n : n \in \mathbb{N}\}$ 

[A monotone decreasing sequence that is bounded below also converges!]

Example:

$$a_n = \sum_{k=1}^n \frac{1}{k^2} \le 1 + \sum_{k=2}^n \frac{1}{k(k-1)} = 1 + \frac{n-1}{n} \le 2 \text{ (where the last equality is established by}$$

induction!), so  $a_n \to L$  for some L [Fact:  $L = \pi^2/6$  ...]

The Squeeze Theorem: If  $a_n \leq b_n \leq c_n$  for every n and  $a_n \to L$  and  $c_n \to L$  as  $n \to \infty$ , then  $b_n \to L$  and  $n \to \infty$ .

So, e.g., since  $-1 \le \sin n \le 1$  for all  $n, \frac{1}{n} \sin n \to 0$  as  $n \to \infty$ .

**Subsequences:** A subsequence amounts to choosing <u>some</u> of the terms of a sequence; formally, a subsequence is  $a_{n_k} = a_{g(k)}$  for some strictly monotone increasing function  $g: \mathbb{N} \to \mathbb{N}$ . Subsequences inherit many of the same properties of the original sequence, for example, boundedness and convergence: if  $a_n \to L$  then  $a_{n_k} \to L$  for every subsequence of  $a_n$ . On the other hand,

Every bounded sequence has a <u>monotonic</u> subsequence. [But we cannot determine beforehand whether or not it will be increasing or descreasing.]

So a bounded sequence always has a convergent subsequence!